

AbleC: An Extensible C Framework

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A refresher and introduction.

Extensible languages frameworks

Allow programmers select the features to be used in their programming languages.

- ▶ new syntax (notations)
- ▶ new semantic analyses
(error-checking, optimization opportunities)
- ▶ new optimizations

Why would anyone want to do that?

Programming language features

General purpose features in, e.g., C or Java

- ▶ assignment statements, loops, if-then-else statements
- ▶ functions (perhaps higher-order) and procedures
- ▶ I/O facilities
- ▶ modules
- ▶ data: integer, strings, arrays, structs

Domain-specific features

- ▶ matrix operations (MATLAB)
- ▶ regular expression matching (Perl, Python)
- ▶ statistics functions (R)
- ▶ computational geometry operations (LN)
- ▶ parallel computing (SISAL, NESL, SAC, PQL etc.)

Many similarities, needless differences.

Working with multiple (domain-specific) languages is a headache.

Extensible languages

Allow programmers select the features to be used in their programming languages.

- ▶ new syntax (notations)
- ▶ new semantic analyses
(error-checking, optimization opportunities)
- ▶ new optimizations

Pick a general purpose host language (e.g. ANSI C),
extend with domain-specific features.

`myProgram.xc` $\implies$ `myProgram.c`

Reliably composable language extensions

- ▶ Extension **developers** work independently.
- ▶ Extension **users** use multiple extensions.
 - ▶ not experts in language design
 - ▶ combinations of extensions must “just work”

Challenges for composable language extensions

1. syntax — context free grammars, assoc. regexs
 - ▶ context-aware scanning [GPCE'07]
 - ▶ modular determinism analysis [PLDI'09]
 - ▶ Copper
2. semantics — attribute grammars
 - ▶ attribute grammars with forwarding, collections and higher-order attributes
 - ▶ set union of specification components
 - ▶ sets of productions, non-terminals, attributes
 - ▶ sets of attribute defining equations, on a production
 - ▶ sets of equations contributing values to a single attribute
 - ▶ modular well-definedness analysis [SLE'12a]
 - ▶ modular termination analysis [SLE'12b, Krishnan-PhD]
 - ▶ Silver

Building a parser from composed specifications.

$$\dots \text{CFG}^H \cup^* \{ \text{CFG}^{E_1}, \dots, \text{CFG}^{E_n} \}$$

$$\begin{aligned} & \forall i \in [1, n]. \text{isComposable}(\text{CFG}^H, \text{CFG}^{E_i}) \wedge \\ & \quad \text{conflictFree}(\text{CFG}^H \cup \text{CFG}^{E_i}) \\ \Rightarrow & \quad \text{conflictFree}(\text{CFG}^H \cup \{ \text{CFG}^{E_1}, \dots, \text{CFG}^{E_n} \}) \end{aligned}$$

- ▶ Monolithic analysis - not too hard, but not too useful.
- ▶ Modular analysis - harder, but required [PLDI'09].
- ▶ **Non-commutative** composition of restricted LALR(1) grammars.

Building an attribute grammar evaluator from composed specifications.

$$\dots AG^H \cup^* \{AG^{E_1}, \dots, AG^{E_n}\}$$

$$\begin{aligned} &\forall i \in [1, n]. \text{modComplete}(AG^H, AG^{E_i}) \\ \Rightarrow &\quad \text{complete}(AG^H \cup \{AG_1^E, \dots, AG_n^E\}) \end{aligned}$$

- ▶ Monolithic analysis - not too hard, but not too useful.
- ▶ Modular analysis - harder, but required [SLE'12a].

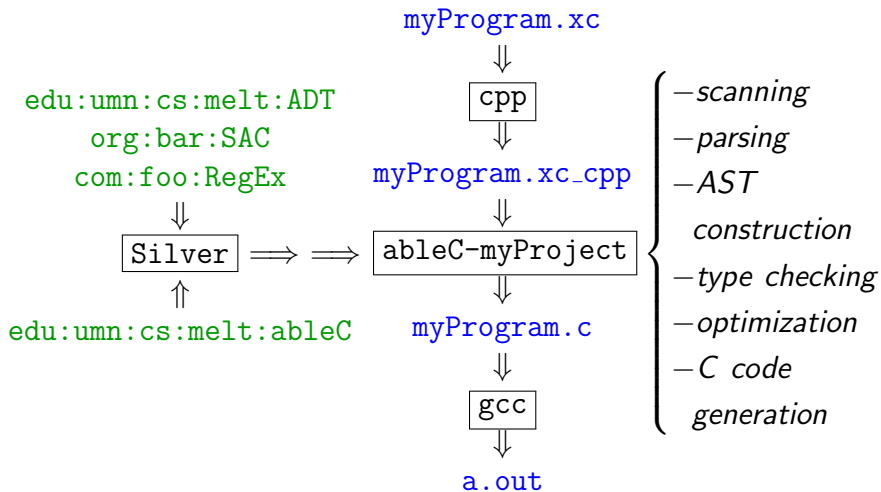
ableC:

Attribute grammar-Based Language Extensions for C

putting these ideas into practice

ableC: a specification of ISO C11 in Silver/Copper.

ableC



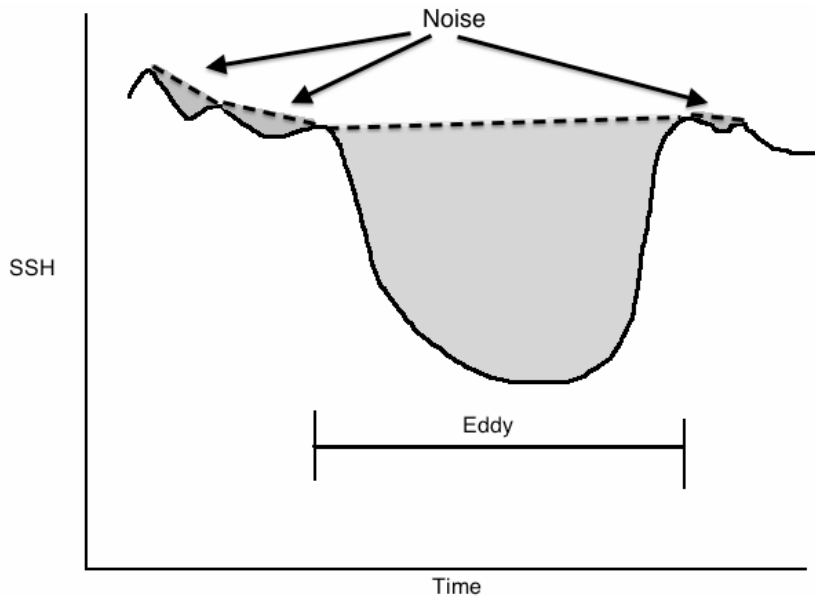
Sample language extensions

An agnostic's view.

Regular expressions

```
1#include "stdio.h"
2#include "regex.h"
3
4int main (int argc, char *argv[]) {
5
6    char *text = readFileContents("X.data") ;
7
8    // eukaryotic messenger RNA sequences
9    regex foo = /^ATG[ATGC]{3,10}A{5,10}$/ ;
10
11    if ( text =~ foo )
12        printf (" Matches...\n" ) ;
13    else
14        printf (" Doesn't match...\n" ) ;
15}
```

A time slice for a point in the ocean



```

1 Matrix float<1> computeArea
2   (Matrix float<1> areaOfInterest)
3 {
4   float y1 = areaOfInterest[0];
5   float y2 = areaOfInterest[end];
6   int x1 = 0;
7   int x2=dimSize(areaOfInterest,0)-1;
8   float m = (y1-y2) / ((float)(x1-x2));
9   float b = y1 - m*x1;
10  Matrix float<1> line = (x1::x2)*m+b;
11  float area
12    = with( x1 <= i < x2)
13      fold(+, 0.0, line - areaOfInterest);
14  return
15    with( 0 <= i < dimSize(line,0) )
16      genarray([dimSize(line, 0)], area);
17 }

```

Halide-inspired transformations

- ▶ Halide: a C++ framework that gives programmers constructs to manipulate nested for-loops for image processing applications.
- ▶ It assumes that the C++ compiler will not generate code with high enough performance.
- ▶ The compiler can perform sophisticated optimizations, but does know now how to optimally apply them.
- ▶ Not for those seeking “casual parallelism”

```
1#include <stdlib.h>
2
3int main() {
4
5    int* grad =
6        (int*) malloc(100*100 * sizeof(int));
7
8    transform grad(x(0->100),y(0->100)) = x + y
9    with { split x by 4, xo, xi;
10           split y by 4, yo, yi;
11           reorder xo, yo, xi, yi
12           }
13
14    // best paired with an extension
15    // for 'proper' matrices
16}
```

```
1#include <stdlib.h>
2
3int main() {
4
5    int* grad =
6        (int*) malloc(100*100*100 * sizeof(int));
7
8    transform grad(x(0->100),y(0->100),z(0->100))
9        = x + y + z
10        with { split x by 4, xo, xi;
11                split y by 4, yo, yi;
12                split z by 4, zo, zi;
13                reorder xo, yo, zo, xi, yi, zi
14            }
15}
```

```

1#include <stdlib.h>
2
3int main() {
4
5    int* grad =
6        (int*) malloc(100*100 * sizeof(int));
7
8    transform grad(x(0->100),y(0->100)) = x + y
9        with { split x by 4, xo, xi;
10               split y by 4, yo, yi;
11               reorder xo, yo, xi, yi;
12               parallelize xo
13           }
14}

```

```
1#include <stdlib.h>
2
3int main() {
4
5    int *grad =
6        (int*) malloc(100*100 * sizeof(int));
7
8    transform grad(x(0->100),y(0->100)) = x + y
9        with { tile x, y, xo, yo, xi, yi, 4, 4 }
10 }
```

Transformation extension

- ▶ These transformations generate the hard-to-write but expected nest of for loops.
- ▶ a reasonable intermediate language for high performance DSL-extensions
- ▶ Halide includes constructs for GPGPU targets, handling ghost cells and storing intermediate results for reuse
all things we could add here

ADTs

- ▶ Implementing algebraic data types in C is sub-optimal.
 - ▶ syntactically verbose
 - ▶ error prone
 - ▶ unsafe
- ▶ A nice opportunity for a language extension.

```
1 /* Simple expresssion evaluation
2    using algebraic datatype as in
3    ML or Haskell. */
4
5 #include <stdio.h>
6 #include <stdlib.h>
7
8 typedef  datatype Expr  Expr;
9
10 datatype Expr {
11     Add (Expr*, Expr*);
12     Sub (Expr*, Expr*);
13     Mul (Expr*, Expr*);
14     Div (Expr*, Expr*);
15     Const (int);
16 };
17
```

```
18 int value (Expr *e) {
19     int res;
20     match (e) {
21         Add(e1, e2): res = value(e1) + value(e2);
22         Sub(e1, e2): res = value(e1) - value(e2);
23         Mul(e1, e2): res = value(e1) * value(e2);
24         Div(e1, e2): res = value(e1) / value(e2);
25         Const(v):   res = v;
26     };
27     return res;
28 }
29 int main () {
30     Expr *t = Add( Const(3),
31                   Mul(Const(2), Const(4)) ) ;
32     printf("value is %d\n", value(t) );
33     return 0;
34 }
```

ADT extension

- ▶ This extension generates the expected collection of
 - ▶ `struct` and `union` data structures.
 - ▶ constructor functions
- ▶ Coming soon
 - ▶ semantic analysis
 - ▶ disallowing access to underlying representation
 - ▶ pattern matching completeness checking

Questions.

ableC: a platform for DSL-feature implementation and distribution.

- ▶ What extensions should we build?
- ▶ What extensions might you like to build?
We aren't domain experts in all of these areas!
- ▶ Are the restrictions too onerous?
- ▶ What can't you do?
 - ▶ language “modifications”
 - ▶ e.g. adding garbage collection for existing pointers

Questions?

Thanks for your attention.

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