

Advanced Stencil-Code Generation















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The Challenges of Exascale Computing

Exascale

- Goal for 2018–2020: 10¹⁸ FLOP/s
- Present fastest supercomputers: 10¹⁶ FLOPS/s

Special Challenges

- massive parallelism on-chip: hundreds of cores
- millions of cores in the cluster
- heterogeneous architectures: different types of cores
- power conservation
- big data management
- robust and performance-portable software



Two Alternative Approaches in SPPEXA

- The Conservative Approach
 - starting point: MPI, OpenMP and variants, OpenCL, Pthreads, ...
 - develop the above incrementally towards exascale by tuning their implementations adding some suitable features

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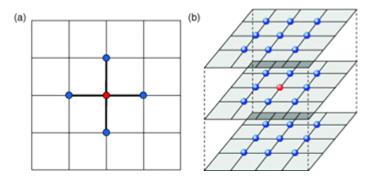
The Radical Approach

- more various, abstract and domain-specific programming languages
- powerful optimizing translation technology for small domains
- more flexible and intelligent run-time systems
- flexibly constructable and usable software development tools
- more performance-portable implementations of application software



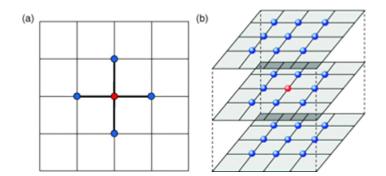
Project Charter

- Domain: Stencil codes
 - widely used
 - well delineated
 - massively parallel



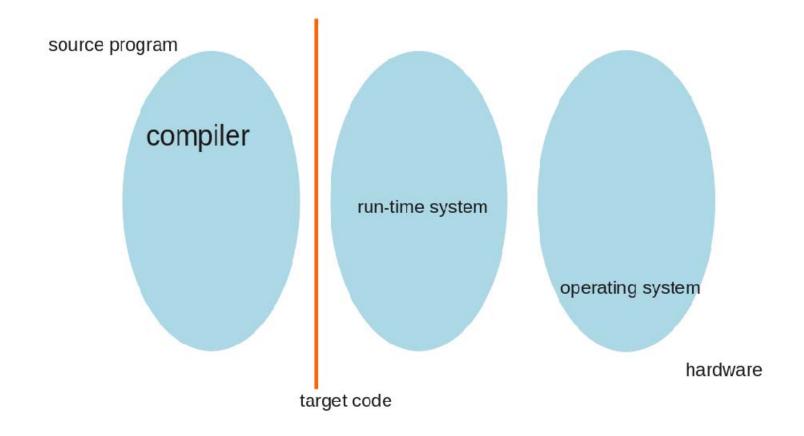
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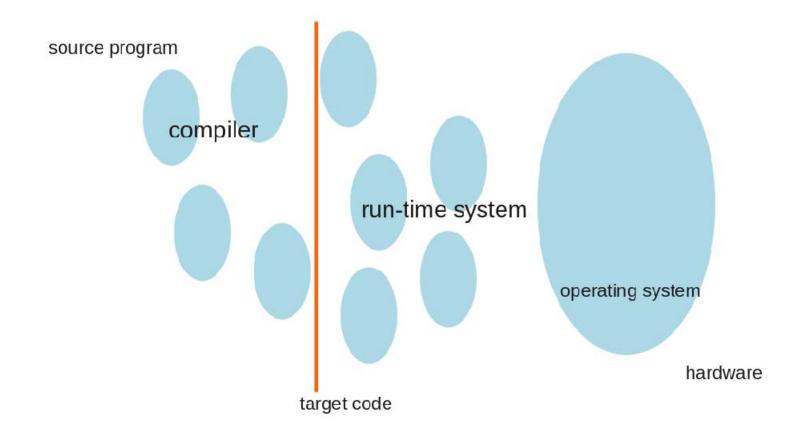


- Approach: Co-design of application, algorithm and software
 - domain-specific optimization at every step of the design
 - use of a variety of languages at various levels of abstraction
 - high automation via customized tool support
 - inspiration: project SPIRAL (optimizing compiler for linear transforms)

The Past: Power to the Compiler



The Future: More Power to the Run-time System





A: Algorithmic engineering

B: Domain-specific representation and modeling

C: Domain-specific optimization and generation

D: Polyhedral optimization and code generation

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E: Platform-specific code optimization and generation

Choice of PDE operators

Choice and tuning of solvers

Focus on multigrid

Identification of variants (features)
 algorithm
 architecture

$$\phi_{\text{MGM}}^{(k)}(\mathbf{x}_k, \mathbf{b}_k) = (\phi_{\text{S}}^{(k)})^{\nu_2} ((\phi_{\text{S}}^{(k)})^{\nu_1}(\mathbf{x}_k, \mathbf{b}_k) + P_k ((\phi_{\text{MGM}}^{(k-1)})^{\gamma}(\mathbf{0}, R_k(\mathbf{b}_k - A_k\mathbf{x}_k))), \mathbf{b}_k)$$

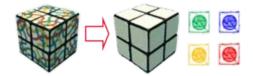
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- B: Domain-specific representation and modeling
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- D: Polyhedral optimization and code generation
- E: Platform-specific code optimization and generation
- Most abstract executable representation algorithmic features architectural features
- Application view

A: Algorithmic engineering

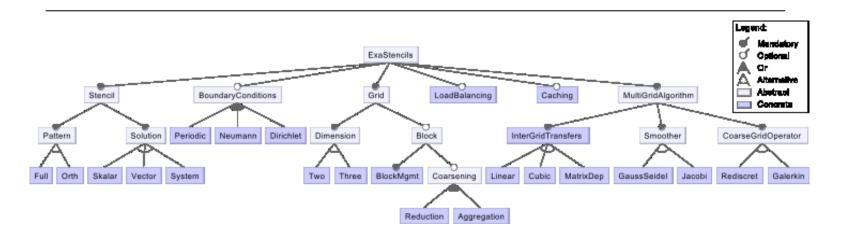
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- Feature orientation
- Exploitation of program similarity
- Product-specific optimization
- Automatic product weaving



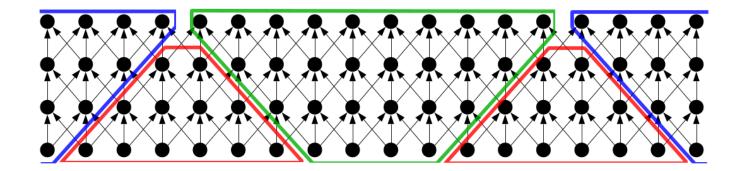
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- Product-specific optimization
- Platform-specific optimization intra-node inter-node



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- Addresses heterogeneity
- Methods considered:

Autotuning Machine learning

Hardware/software codesign

```
void GaussSeidel_rb(int lev, Array <double> *Sol, Array <double> *RHS) {
   int offset;

#pragma omp parallel for private(offset)
   for (int i=1; i < Sol[lev].nrows()-1; i++) {
      offset = (i % 2 == 0 ? 2 : 1);
      for (int j = offset; j < Sol[lev].ncols()-1; j+=2) {
            Sol[lev] (i, j) = double(0.25)*(RHS[lev] (i, j) + Sol[lev] (i+1, j) + Sol[lev] (i-1, j) + Sol[lev] (i, j+1) + Sol[lev] (i, j-1));
      }
}

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      }
}
}</pre>
```

Supercomputers

SuperMUC

- Leibniz Computation Centre
- 6th on TOP500, 11/2012



- Jülich Research Centre
- 5th on TOP500, 11/2012



- Tokyo Institute of Technology
- 17th on TOP500, 11/2012

