

Programming Languages through the Lens of Program Generation

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Quicksort – in C

```
void quickSort( int a[], int l, int r) {
    int j;
    if( l < r ) {
        j = partition( a, l, r);
        quickSort( a, l, j-1);
        quickSort( a, j+1, r);
    }
}

int partition( int a[], int l, int r) {
    int pivot, i, j, t;
    pivot = a[l];
    i = l; j = r+1;
    while( 1 ) {
        do ++i; while( a[i] <= pivot && i <= r );
        do --j; while( a[j] > pivot );
        if( i >= j ) break;
        t = a[i]; a[i] = a[j]; a[j] = t;
    }
    t = a[l]; a[l] = a[j]; a[j] = t;
    return j;
}
```

Quicksort – in Haskell

```
import Data.List (partition)

quicksort [] = []
quicksort (x:xs) = quicksort less ++ (x:quicksort rest)
    where (less, rest) = partition (< x) xs
```

Quicksort – in λ -Prolog

```
quicksort([], []).
```

```
quicksort([H | T], S) :-  
  partition(>(H), T, L, R),  
  quicksort(L, SL),  
  quicksort(R, SR),  
  append(SL, [H | SR], S).
```

Quicksort – in Scheme

```
(define pivot (lambda (l)
  (cond ((null? l) 'done)
        ((null? (cdr l)) 'done)
        ((<= (car l) (cadr l)) (pivot (cdr l)))
        (#t (car l)))))

(define partition (lambda (piv l p1 p2)
  (if (null? l) (list p1 p2)
      (if (< (car l) piv) (partition piv (cdr l) (cons (car l) p1) p2)
          (partition piv (cdr l) p1 (cons (car l) p2))))))

(define (quicksort l)
  (let ((piv (pivot l)))
    (if (equal? piv 'done) l
        (let ((parts (partition piv l () ())))
          (append (quicksort (car parts))
                  (quicksort (cadr parts)))))))
```

Quicksort – in Scala I

```
class QuickSort[A <% Ordered[A]] extends SortingAlgorithm[A] {  
  def sort(l:List[A]): List[A] = {  
    if(l.size <= 1) return l  
    val pivot = l.head  
    val (lt,gt) = l.tail.partition( _ <= pivot)  
    sort(lt) ::: List(pivot) ::: sort(gt)  
  }  
}
```

Quicksort – in Scala II

```
class QuickSort[A <% Ordered[A]] {  
  def sort(xs: Array[A]): Array[A] {  
    def swap(i: Int, j: Int) {  
      val t = xs(i); xs(i) = xs(j); xs(j) = t  
    }  
    def sort1(l: Int, r: Int) {  
      val pivot = xs((l + r) / 2) var i = l; var j = r  
      while (i <= j) {  
        while (xs(i) < pivot) i += 1  
        while (xs(j) > pivot) j -= 1  
        if (i <= j) {  
          swap(i, j)  
          i += 1  
          j -= 1  
        }  
      }  
      if (l < j) sort1(l, j)  
      if (j < r) sort1(i, r)  
    }  
    sort1(0, xs.length - 1)  
  }  
}
```

Quicksort – in Haskell II

```
partition arr left right pivotIndex = do
    pivotValue <- readArray arr pivotIndex
    swap arr pivotIndex right
    storeIndex <- foreachWith [left..right-1] left (\i storeIndex -> do
        val <- readArray arr i
        if (val <= pivotValue)
            then do
                swap arr i storeIndex
                return (storeIndex + 1)
            else do return storeIndex )
    swap arr storeIndex right
    return storeIndex

qsort arr left right = when (right > left) $ do
    let pivotIndex = left + ((right-left) `div` 2)
    newPivot <- partition arr left right pivotIndex
    qsort arr left (newPivot - 1)
    qsort arr (newPivot + 1) right

swap arr left right = do
    leftVal <- readArray arr left
    rightVal <- readArray arr right
    writeArray arr left rightVal
    writeArray arr right leftVal
```

And others too

- ▶ Hello World
- ▶ max/min
- ▶ powering
- ▶ fibonacci
- ▶ factorial
- ▶ gcd/lcm
- ▶ palindrome
- ▶ Gaussian Elimination
- ▶ λ -calculus interpreter

Will the real Quicksort please stand?

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1. Data traversal
2. Divide and conquer strategy
3. End result (a sorted container of data)

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► Different:

1. Language
2. Data Structure
3. Pivot choice (first, middle, ...)
4. Partition: **In-place** or Copy
5. Partition: recursive or iterative
6. Use of library functions
7. Representation of current 'focus' (implicit/explicit)

(and the differences are not independent)

What we have

A ((pretty) source) code generator,

- ▶ written in Haskell,
- ▶ for: quicksort (plus partition, swap, concat), Hello World, max/min, powering, fibonacci, factorial, gcd/lcm, palindrome.
- ▶ producing: C, C[‡], Java, Lua, Lisp, Haskell, LambdaProlog,
- ▶ parametrized by **design decisions**.

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Where the “program” input is roughly

```
quicksort(items)
  if items.length < 1 return items
  choose index
  pivot := items[index]
  partition pivot less rest
  return quicksort(less) ++ (pivot ++ quicksort(rest))
```

Full disclosure: current Haskell output

Choices: Pivot = head, DS = list, lib = [Part, Concat]

```
module QuicksortProg where
import Data.List (partition)
quicksort items len =
  if (<) len 1
  then items
  else quicksortRec items 0 (len-1)
quicksortRec items leftIndex rightIndex =
  if (>=) leftIndex rightIndex
  then items
  else
    let pivot = (head items) in
    let (less,rest) = partition (< pivot) (tail items) in
    quicksortRec less 0 ((-) (length less) 1)
      ++ (pivot:quicksortRec rest 0 ((-) (length rest) 1))
```

Generalizing

program family = concepts + design choices + staging

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Concepts present in implementation and theory of quicksort:

- ▶ (indexable) container
- ▶ ordering / comparison
- ▶ partition, permutation, pivot
- ▶ choice (of an index)
- ▶ data traversal
- ▶ memory access (for both read and write)
- ▶ (baked in) divide and conquer strategy

Generalizing

program family = concepts + design choices + staging

Design choices:

- ▶ Implementation Language
- ▶ Data Structure (list ,array)
- ▶ Pivot choice (first, middle, ...)
- ▶ Partition: recursive or iterative
- ▶ Use of library functions for *partition* and/or *concat*

Generalizing

program family = concepts + design choices + staging

Staging

- ▶ Hard to describe – no good theory of heterogeneous staging
- ▶ 5 “stages”: meta-theory, concept, design, code, run-time.
- ▶ Most important part - the **weaving instructions**:

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What's wrong with current PLs?

Some are too low level

- ▶ must worry about memory use and traversal
- ▶ cannot write generic code
- ▶ rigid syntax
- ▶ no access to (convenient) high-level concepts

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- ▶ so much overloading that code is impossible to read
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Opinion: largely caused by:

- ▶ no proper understanding of **staging**
- ▶ too high a focus on **general purpose**

Opportunities

Your Language is my Assembler

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Concept language and library

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Proper language of Algorithms