A Formalization of Moessner's Theorem in Coq

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The stream of odd natural numbers:

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The corresponding stream of partial sums:

1,

The stream of odd natural numbers:

The corresponding stream of partial sums:

i.e.,

$$1^2$$
, 2^2 , 3^2 , 4^2 , 5^2 , 6^2 , 7^2 , 8^2 , ...

Constructively

- start from the stream of natural numbers
- strike out every 2nd element
- compute the successive partial sums

Result: the stream of squares .

Scaling up

- start from the stream of natural numbers
- strike out every 3rd element
- compute the successive partial sums
- strike out every 2nd element
- compute the successive partial sums

Result: the stream of ...

Scaling up

- start from the stream of natural numbers
- strike out every 3rd element
- compute the successive partial sums
- strike out every 2nd element
- compute the successive partial sums

Result: the stream of cubes .

Scaling up: Moessner's theorem

- start from the stream of natural numbers
- strike out every nth element & sum
- ullet strike out every $(n-1){\rm th}$ element & sum
- ...
- strike out every 3rd element & sum
- strike out every 2nd element & sum

Result: the stream of powers of n.

Background

- Moessner (1951)
 The property, no proofs.
- Perron (1951), Paasche (1952), Salié (1952)
 Complicated inductive proofs.
- Hinze (IFL 2008)
 A calculational proof.
- Rutten & Niqui (HOSC 2012)
 A co-inductive proof.

This work (in progress)

- a formalization in Coq
- but first, learning Coq

Learning Coq in principle

- web resources
- book
- seasonal schools

Learning Coq in practice

- practice, practice, practice
- need a TA (or ideally, a coach)
- forces you to think things through
- can be frustrating at times
- wonderfully rewarding, overall

Teaching Coq

- introduction to functional programming (Q3)
- more advanced functional programming (Q4)

a marvelous experience

Term projects in Q3

- a standard batch (interpreters, compilers, decompilers, VMs, CPS, power series, searching in binary trees, Boolean negational normalization, FSA, etc.)
- a cherry on top of the pie: formalizing a theorem and a proof from another course(!)

Term projects in Q4

- functional & relational programming
- the Ackermann-Peter function
- Boolean normalization and equisatisfiability
- abstract interpretation (strided intervals)
- group theory and pronic numbers
- B-trees
- graph theory
- reduction from circuit to SAT
- vector spaces & Cauchy-Schwarz inequality

Plan

- Moessner's theorem at degree 3
- Moessner's theorem at degree 4
- Which starting indices in the master lemma?
- Conclusion and perspectives

We are given

- stream_of_ones
- stream_of_positive_powers_of_3
- stream_of_positive_powers_of_4
- skip_2, skip_3, skip_4, ...
- sums & sums_aux, which uses an accumulator
- stream_bisimilar

Moessner's theorem at degree 3

- the statement
- the proof
- the master lemma

```
Proof.
  unfold stream_of_positive_powers_of_3.
  unfold sums.
  apply (Moessner_3_aux 0).
Qed.
```

```
Lemma Moessner_3_aux :
  forall (n : nat),
    stream_bisimilar
      (make_stream_of_nats (S n)
                        (fun i => i * i * i)
                            S)
      (sums_aux ???
        (skip_2
          (sums_aux ???
            (skip_3
               (sums_aux ???
                stream_of_ones))))).
```

Moessner's theorem at degree 4

- the statement
- the proof
- the master lemma

```
Theorem Moessner_4:
  stream_bisimilar
    stream_of_positive_powers_of_4
    (sums
      (skip_2
        (sums
          (skip_3
            (sums
              (skip_4
                 (sums
                   stream_of_ones)))))).
```

```
Proof.
  unfold stream_of_positive_powers_of_4.
  unfold sums.
  apply (Moessner_4_aux 0).
Qed.
```

```
Lemma Moessner_4_aux :
  forall (n : nat),
    stream_bisimilar
      (make_stream_of_nats (S n)
                    (fun i => i * i * i * i)
                            S)
      (sums_aux ???
         (skip_2
          (sums_aux ???
            (skip_3
               (sums_aux ???
                 (skip_4
                   (sums_aux ???
                     stream_of_ones)))))).
```

So, which starting indices?

Reminder: Newton's binomial expansion

$$(n+1)^{2} = n^{2} + 2 \cdot n + 1$$

$$(n+1)^{3} = n^{3} + 3 \cdot n^{2} + 3 \cdot n + 1$$

$$(n+1)^{4} = n^{4} + 4 \cdot n^{3} + 6 \cdot n^{2} + 4 \cdot n + 1$$

```
Lemma Moessner_2_aux :
  forall (n : nat),
    stream_bisimilar
      (make_stream_of_nats (S n)
                              (fun i \Rightarrow i * i)
                             S)
      (sums_aux (n * n))
         (skip_2
           (sums_aux (2 * n))
             stream_of_ ones ))).
```

```
Lemma Moessner_3_aux :
  forall (n : nat),
    stream_bisimilar
      (make_stream_of_nats (S n)
                        (fun i => i * i * i)
                            S)
      (sums_aux (n * n * n)
        (skip_2
          (sums_aux (3 * n * n))
            (skip_3
               (sums_aux (3 * n))
                stream_of_ ones ))))).
```

```
Lemma Moessner_4_aux :
  forall (n : nat),
    stream_bisimilar
      (make_stream_of_nats (S n)
                    (fun i => i * i * i * i)
                            S)
      (sums_aux (n * n * n * n))
         (skip_2
          (sums_aux (4 * n * n * n))
            (skip_3
               (sums_aux (6 * n * n))
                 (skip_4
                   (sums_aux (4 * n))
                     stream_of_ ones ))))))).
```

Assessment

- The monomials of the binomial expansion.
- At this point, this is still a conjecture.
- Essential algebraic support from Coq to find the starting indices.

Conclusion and perspectives

Proof assistants are changing the world:

- the 4-color theorem
- G. Gonthier's work at Microsoft Research
- C. Schuermann's DemTech in Denmark
- V. Voevodsky's work at Princeton

This was then: MFPS 1991

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"Look guys.

We don't say anything about your proofs, so don't say anything about our programs, OK?"

Moi: A lightweight question: how does it feel, ...

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Thank you.